

# Unbiased generation of metastable states for Ising spin glasses

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# Overview

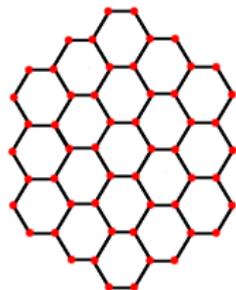
- ▶ Background and motivation
- ▶ Basics of spin glass lattices
- ▶ Equivalence between edge coloring and spin configuration
- ▶ Generating the metastable states
- ▶ Generalization idea
- ▶ Estimating the number of local minima
- ▶ Knuth's method vs. exact calculations
- ▶ Ferromagnet - spin glass phase transition
- ▶ Sampling the energy distribution
- ▶ Summary and future work

## Motivation, background

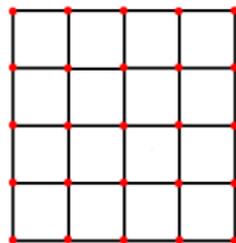
- ▶ An efficient way to generate all metastable states of a lattice
- ▶ Knuth's algorithm for estimating the size of a search tree
- ▶ Phase transition between ferromagnet and spin glass
- ▶ Is the number of metastable states affected?
- ▶ Various applications:
  - ▶ Estimating the number of minima as a function of system size
  - ▶ Exploring the energy distribution of minima and sampling

## Models

We consider the following two types of 2D Ising spin glass lattices:



(a) Hex lattice

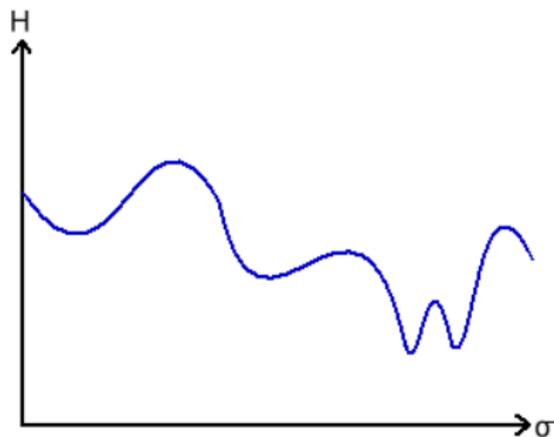


(b) Square lattice

- ▶ Up-down spins:  $\sigma_i \in \{1, -1\}$
- ▶ Positive-negative interactions:  $J_{ij} \in \{1, -1, 0\}$
- ▶ Hamiltonian:  $H(\sigma) = -\frac{1}{2} \sum_{ij}^N J_{ij} \sigma_i \sigma_j$

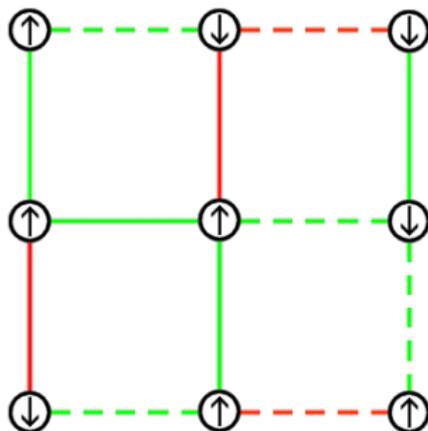
# Basics

- ▶ Local minimum: no energy decrease by a single spin flip
- ▶ Usually wanted: ground state,  $\min(H(\sigma))$ 
  - ▶ Hard
- ▶ Greedy algorithms: get stuck in local minima
- ▶ Ground state(s) hard to prove



## Graphical view of local minima (1/2)

- ▶ A spin configuration  $\sigma$  induces a coloring of the edges in the respective lattice into “green” (satisfied) and “red” (unsatisfied)

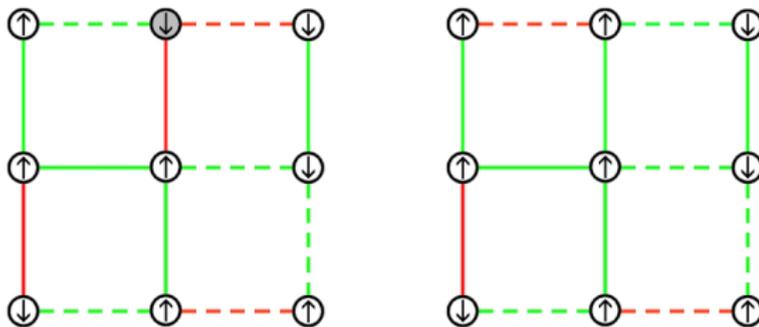


Here solid edge indicates a positive interaction, a dotted edge a negative one.

- ▶ *Note:*  $H = -(\# \text{ of green edges}) + (\# \text{ of red edges})$

## Graphical view of local minima (2/2)

- ▶ An edge coloring of the lattice corresponds to a (proper) local energy minimum, if and only if each vertex is incident to (properly) fewer red than green edges.
- ▶ E.g. the previous example coloring (spin configuration) can still improved:



- ▶ However, all colorings are not valid
- ▶ The sum of negative interactions and red colorings must be even number for each cycle

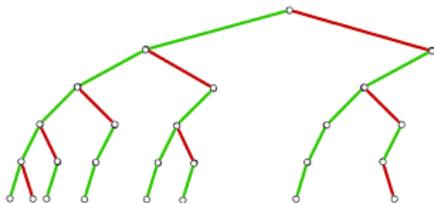
## Generating local minima (1/2)

- ▶ A local minimum corresponds to an edge coloring of the lattice (interaction graph)  $G$  satisfying a local consistency condition.
- ▶ Two minima per each valid coloring
- ▶ Consider any spanning tree  $T$  of  $G$ .
- ▶ Any edge colouring of  $G$  obviously yields a unique colouring of  $T$  and vice versa.
- ▶ Finding a local minimum of  $T$  is easy, but the corresponding coloring of  $G$  is not necessarily a local minimum.

## Generating local minima (2/2)

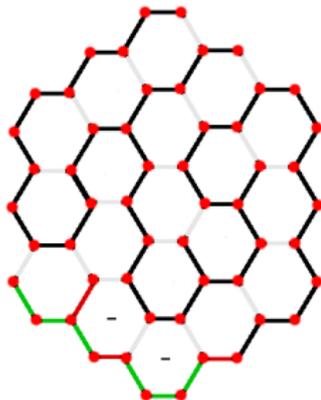
An algorithm to generate the local minima on an interaction graph  $G$ :

- ▶ Choose spanning tree  $T$  of  $G$ . List edges of  $T$  in some order.
- ▶ Enumerate systematically all colourings of this edge list ( $\implies$  binary search tree).
- ▶ Only consider branches of the search tree where the current partial colouring can be completed to a valid colouring of all of  $G$ . (This may be difficult to test!)
- ▶ Now leaves of the search tree correspond one-to-one to consistent colourings ( $\sim$  local minima) of  $G$ .



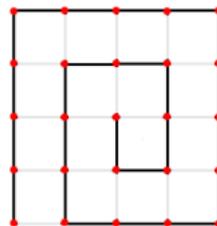
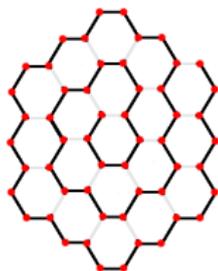
## Complexity of completability

- ▶ In a general spanning tree, determining the completability of a given partial colouring may require arbitrary long look-ahead:



## Spanning paths

- ▶ In lattices, spiral-like paths provide particularly simple spanning trees:



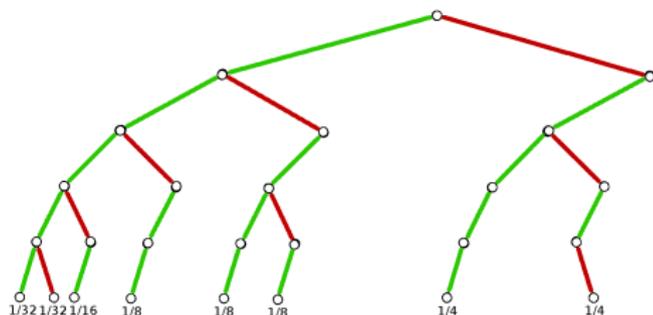
- ▶ Moreover, if the edges of the spanning spiral are coloured inside-out, the completability can be tested efficiently. (Basically, a partial colouring can always be completed, unless the validity conditions are violated by already chosen or “locally forced” edge colours.)

## Generalization of the tree generation

- ▶ Enumerate the edges of  $G$
- ▶ For each node  $v$  generate a SAT-formula “at most  $d(v)/2$  incident edges may be unsatisfied”
- ▶ For each uncolored edge:
  - ▶ Try green and red colorings
  - ▶ Apply propagation rules
    - ▶ SAT-formulas are basically implications
    - ▶ Validity of coloring
  - ▶ If a contradiction follows, prune the branch
  - ▶ Randomly select an available branch
- ▶ Coming up with a proper set of propagation rules for arbitrary graph is still work in progress
- ▶ The idea is to detect the dead-ends in advance so backtracking would not be necessary

# Estimating the size of a search tree

- ▶ D. E. Knuth, “Estimating the efficiency of backtrack programs”. Math. Computation 29 (1975), 121–136.
- ▶ Method to estimate number of leaves  $S$  in a large search tree  $T$ :
  - ▶ Make a random descent into  $T$ , starting from the root.
  - ▶ Record degrees (branching factors)  $d_1, d_2, \dots, d_n$  of the vertices encountered along the descent. (For a binary tree,  $d_i \in \{1, 2\}$  for each  $i$ .)
  - ▶ Compute estimate  $\hat{S} = d_1 d_2 \cdots d_n$ .
- ▶ *Theorem.*  $\hat{S}$  is an unbiased estimate of the true tree size  $S$ .



## Reliability of the Knuth's method

- ▶ How many descents do we need?
- ▶ Different ways to decrease the variance:
  - ▶ Increase the number of descents
  - ▶ Modify the search tree (using lattice structure information)
  - ▶ Perform biased descents (if we know that the search tree is biased)
- ▶ For better understanding, we performed some exact calculations and compared results

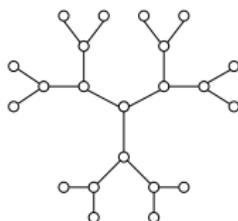
## Exact calculations (1/2)

Ferromagnetical  $L \times L$  square lattice

- ▶ Local minima form horizontal or vertical red “ladders” through the lattice
- ▶ Let  $x_L$  be the number of ways to put vertical ladders through a  $L \times L$  lattice
- ▶ Recursion:  $x_L = x_{L-1} + x_{L-2} + 1$ ,  $x_0 = 0$ ,  $x_1 = 1$
- ▶ Interestingly,  $x$  happens to be similar to Fibonacci series  $f$
- ▶  $x = \{0, 1, 2, 4, 7, 12, 20, 33, \dots\}$
- ▶  $f = \{0, 1, 1, 2, 3, 5, 8, 13, 21, 34, \dots\}$
- ▶  $f_i = x_{i-2} + 1$
- ▶ Number of local minima is  $2(2x_L + 1)$  (vertical or horizontal ladders, the ground state and up-down symmetry)
- ▶  $\frac{4}{\sqrt{5}}(\phi^{L+2} - (1 - \phi)^{L+2}) - 2$ ,  $\phi = (1 + \sqrt{5})/2$

# Exact calculations (2/2)

Bethe lattice ( $d = 3$ )

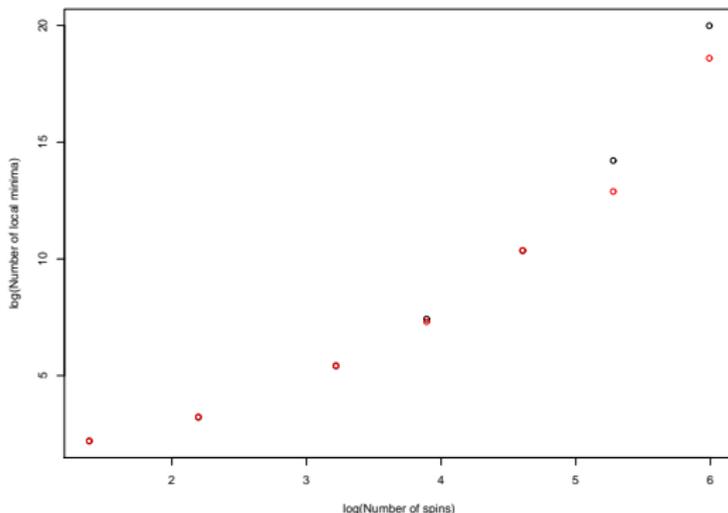


- ▶ Number of local minima can again be solved by recursion
- ▶  $y_n = y_{n-1}^2 + 2y_{n-1}y_{n-2}^2$
- ▶  $x_n = y_{n-1}^3 + 3y_{n-1}^2y_{n-2}^2$
- ▶  $y_1 = 1, y_2 = 3$
- ▶ Number of local minima for a Bethe lattice with radius  $i$  is then  $2x_i$  (up-down symmetry)
- ▶  $2x = \{2, 8, 108, 18900, 5.73358 \cdot 10^8, \dots\}$

# Comparing the results (1/2)

Ferromagnetical  $L \times L$  square lattice (10000 descents per lattice)

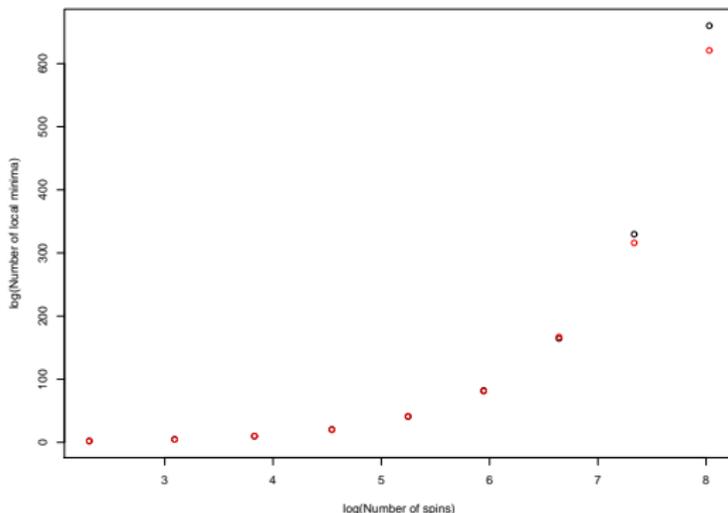
- ▶ Search tree is very biased
- ▶ In the following graph black dots are the real number of local minima and red dots are the estimates given by the algorithm:



## Comparing the results (2/2)

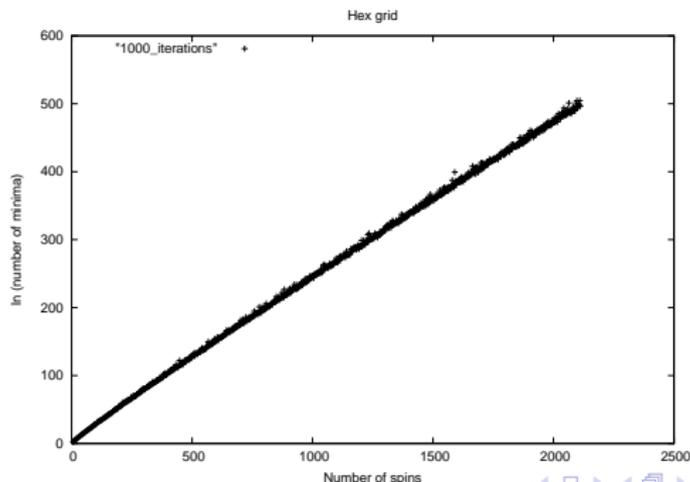
Bethe lattice ( $d = 3$ ) (10000 descents per lattice)

- ▶ Search tree is quite balanced
- ▶ In the following graph black dots are the real number of local minima and red dots are the estimates given by the algorithm:



# Application 1: Estimating the number of minima (1/2)

- ▶ Ansatz:  $S \sim e^{\alpha N}$ , where  $S$  is the number of local minima,  $N$  is the number of spins, and  $\alpha$  is a coefficient which depends on the system characteristics.
- ▶ E.g. for hexagonal lattices of up to 1000 hexes ( $> 2000$  spins), with 50% fraction of negative interactions, we obtain the following diagram. (Each estimate of  $L$  is based on 1000 descents.)

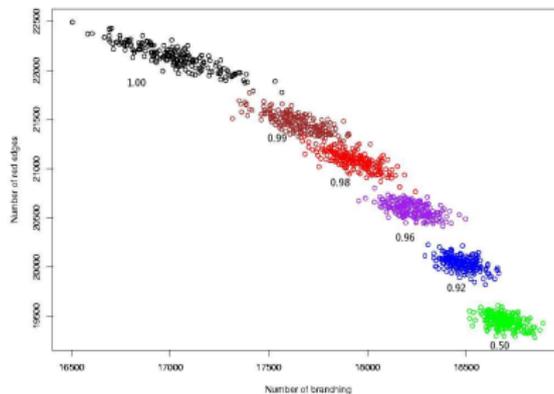


## Application 1: Estimating the number of minima (2/2)

- ▶ Numerical estimates of the coefficient  $\alpha$ , from systems of 1...1000 hexes:
  - ▶ Ferromagnetic hex lattice:  $\alpha \approx 0.226$
  - ▶ Hex lattice with 50% negative interactions:  $\alpha \approx 0.231$
- ▶ The results match quite well the analytic prediction of  $\alpha \approx 0.225$  from a recent paper by Waclaw and Burda (arXiv Jan 2008). The paper also reports experimental data on systems of up to 24 spins, indicating  $\alpha \approx 0.226$ .

## Application 2: Studying the phase transition

- ▶ When a certain fraction of bonds are antiferromagnetical (negative edges) the system will undergo a phase transition.
- ▶ For different systems the fraction of bonds differ
- ▶ Will the number of local minima be affected among other properties?
- ▶ Results from 60000 spin hexagonal lattice:



## Uniform sampling of leaves in a search tree

- ▶ Apply Knuth's method recursively.
- ▶ When at a non-leaf vertex of a binary search tree  $T$ :
  - ▶ Estimate size of left subtree:  $\hat{S}_L$
  - ▶ Estimate size of right subtree:  $\hat{S}_R$
  - ▶ Descend to left subtree with probability

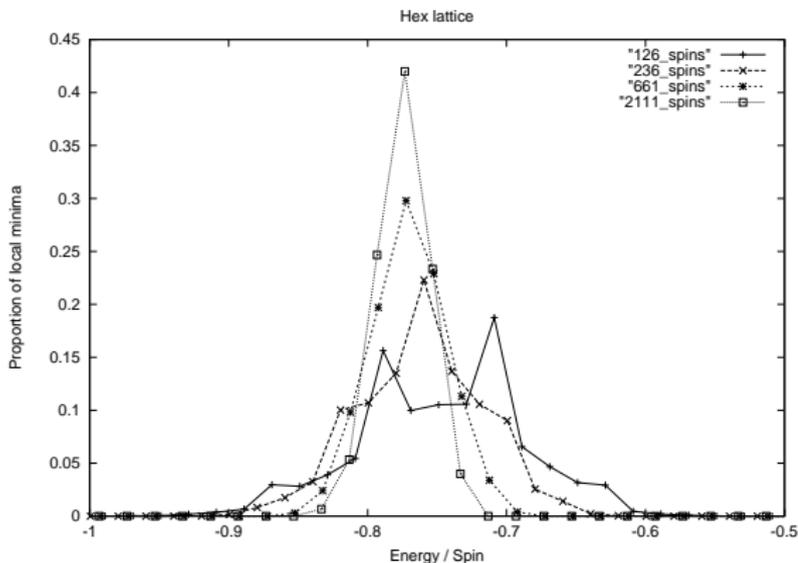
$$p = \frac{\hat{S}_L}{\hat{S}_L + \hat{S}_R},$$

and to right subtree with probability  $1 - p$ .

⇒ Descent reaches each leaf with equal probability.

## Application 3: Energy distribution of local minima

- Based on a uniform sampling of 1000 minima in ferromagnetic hex lattices with  $N = 126, \dots, 2111$  spins. (Equidistant binning of energy levels. Only one descent per subtree size estimate.)



## Summary and future work

- ▶ New combinatorial method for almost uniform sampling of local minima in spin glass lattices.
- ▶ Method has potentially many applications to the exploration of spin glass energy landscapes.
- ▶ Basic method scales well to large system sizes, but search tree shape affects the variance quite much.
- ▶ Future work 1: Better analysis of the sampling and perhaps biased sampling for the full energy distribution.
- ▶ Future work 2: Extend the method from 2D lattices to other graph structures.
- ▶ Future work 3: Explore other applications in landscape analysis.